Dismantlings and Iterated Clique Graphs

Martín Eduardo Frías Armenta * Victor Neumann Lara †

December 5^{th} , 2001

Abstract

The clique graph k(G) of a graph G is the intersection graph of the family of cliques of G. A vertex z is dominated by a neighbor w if every neighbor of z is also neighbor of w. We prove in this paper that deletion of dominated points of G does not change the k-behavior of G.

Keywords: Dominated, invariance, k-behavior.

2000 Mathematics Subject Clasification: 05C99

1 Introduction.

A clique of a graph G is a maximal complete subgraph of G. The clique graph k(G) of a graph G is the intersection graph of the family of cliquues of G. We say that a graph F is k-periodic with period p if is it isomorphic to $k^p(F)$ but not to $k^q(F)$ for $1 \le q < p$. It is said that a graph G is k-stationary if there is an integer n such that $k^n(G)$ is k-periodic. A special case of k-stationary graphs is a k-null graph, G is k-null if there is n such that $k^n(G) \cong K_1$. We say that G k-diverges if $|k^n(G)| \to \infty$ when $n \to \infty$, and in this case we say that $ci(G) = \infty$. The k- behavior of G is determinated by the above clasification i.e., it consists in saying if G is k-stationary, k-null or k-divergent.

Let z, w be elements of graph G. We say that z is dominated by his neighbor w if every neighbor of z is also neighbor of w, that is, $N_G[w] \supseteq N_G[z]$, and we write $w \gtrsim_G z$ (we will omit G when no confusion arisen). We say to that z is dominated if z is dominated by some of its neighbors. Domination induces a preorder. When a graph G do not have dominated points we say that G is irreducible

The present paper was inspired in Prisner [8]. In that paper Prisner proved that if we eliminate dominated points of graph G until to obtain a irreducible

^{*}Universidad de Sonora, Convenio de Retención Conacyt Ref: 489100-1 Exp: 010260

[†]Instituto de Matemáticas, UNAM

graph H, and H have no triangles, then there are n such that $k^n(G) \in \{h, k(H)\}$ and Prisner shows too how to have a bound of n. He use a complicates technics like homology to prove this result. We have found an easier proof of this result whit the technics in present paper, this proof is in [2, 3].

We can resume the main result in this paper as it follows: If $x \in V(G)$ is dominated then G and $G - \{x\}$ have the same k- behavior.

The k-behavior is studied for many reasons, one of these, is the result of Hazan and Neumann-Lara [5]. They establish that every endomorphism in a comparison graph that is k-null has the property of fix point.

For notation see [4] for large bibliography see [2].

2 Dismantlings.

Definition 2.1 Let G be a graph and H' a subgraph of G.

- 1. There is an innerly short dismantling of G to H' if every element of V(G)-V(H') is dominated by some element of V(H'); and we write $G \stackrel{\#_0}{\longrightarrow} H'$
- If H is a graph, we say that there is a short dismantling of G to H when there exist H', subgraph of G, which is isomorphic to H and such that G ^{#0}→ H'. In this case we write G [#]→ H, in particular G [#]→ H'.

We could observe that $c(G) \stackrel{\#_0}{\longrightarrow} k(G)$, where c(G) is the intersecction graph of family of complete subgraphs of G.

Let us observe that G is a cone if and only if $G \stackrel{\#}{\to} K_1$.

3 Dismantlings and the Clique Operator

Lemma 3.1 Assume $G \stackrel{\#_0}{\to} H$ and A, B are cliques of G (not necessarily different) with $A \cap B \neq \emptyset$. Then there is $y \in V(H)$ such that $y \in A \cap B$.

Theorem 3.2 Let G and H be graphs such that $G \stackrel{\#}{\to} H$. Then $k(G) \stackrel{\#}{\to} k(H)$.

Let $H'\cong H$ such that $G\stackrel{\#_0}{\to} H'$. Let us define $g:k(H')\to k(G)$, let $B\in V(k(H'))$. Since B is a complete subgraph of G, we may choose a clique g(B) of G such that $B\subseteq g(B)$. Is easy to see [5,6] that g defines an isomorphism of k(H') in its image.

Now we will prove that $k(G) \stackrel{\#_0}{\longrightarrow} g(k(H'))$.

Let $A \in V(k(G)) - V(g(k(H')))$ and B be a clique of H' that contains $A \cap V(H')$. We will prove that $A \lesssim g(B)$:

Let $C \in N[A]$. Then $A \cap C \neq \emptyset$ and by Lemma 3.1 there is $y \in A \cap C \cap V(H')$. Thus $y \in B$, and consequently $y \in g(B)$. Therefore $g(B) \cap C \neq \emptyset$ and $C \in N[g(B)]$.

4 Main Result

Lemma 4.1 Let G, H and H' be graphs. If $G \stackrel{\#}{\to} H$ and $G \stackrel{\#_0}{\to} H' \cong H$ we will define $G_{\#H}$ as the graph of intersection of $A \cap V(H')$ where A is a clique of G. The next staments are hold:

- 1. $c(H') \stackrel{\#_0}{\to} G_{\#H} \stackrel{\#_0}{\to} k(H')$.
- 2. There is a retraction [5, 6] $f: k(G) \to G_{\#H}$ with inverse $g: G_{\#H} \to k(G)$ with the next properties:
 - a. N[g(f(A))] = N[A] for every $A \in k(G)$.
 - b. $k(G) \stackrel{\#_0}{\rightarrow} g(G_{\#H})$.
- 3. $k^2(G) \cong k(G_{\#H})$.
- 4. If H is a periodic graph of period n, then $k^{n-1}(c(H)) \stackrel{\#}{\to} k^n(G) \stackrel{\#}{\to} H$.
- 1. It is clear from $k(H') \subseteq_* G_{\#H} \subseteq_* c(H')$ and $c(H') \stackrel{\#_0}{\rightarrow} k(H')$.
- 2. Let us define $f(A) = A \cap V(H')$. Let B be an element of $V(G_{\#H})$ and g(B) a clique of G that contains B and such that f(g(B)) = B. By Lemma 3.1, f preserves edges and clearly g is an isomorphism over its image. Therefore f is retraction.

Let $A, B \in V(k(G))$ such that f(A) = f(B). By Lemma 3.1, if $C \in V(k(G))$ is a neighbor of A, then it is a neighbor of B too, and since f(g(f(A))) = f(A), we have N[g(f(A))] = N[A]. In particular $N[A] \subseteq N[g(f(A))]$, and hence $k(G) \stackrel{\#_0}{\longrightarrow} g(G_{\#H})$.

- 3. We have $G_{\#H} \cong g(G_{\#H})$, where g is as g in 2. The statment is followed of 2.
- 4. It is followed of 1, 2 and 3.

Lemma 4.2 Let H be a periodic graph of period n then c(H) is k-stationary and there is m such that $k^m(c(H)) \cong k^{m+n}(c(H))$.

We have that $c(H) \stackrel{\#}{\to} k(H)$. In Lemma 4.1.4 we make $G = k^{n-1}(c(H))$ and we obtain

$$k^{n-1}(c(H)) \xrightarrow{\#} k^{2n-1}(c(H)) \xrightarrow{\#} H$$

hence

$$k^{n-1}(c(H)) \xrightarrow{\#} k^{2n-1}(c(H)) \xrightarrow{\#} k^{3n-1}(c(H)) \xrightarrow{\#} \dots \xrightarrow{\#} k^{mn-1}(c(H)) \xrightarrow{\#} H$$

for every natural m, since $k^{n-1}(c(H))$ is of finite order there are two distincts naturals i_0 and i_1 such that:

$$k^{i_0 n - 1}(c(H)) = k^{i_1 n - 1}(c(H))$$

with $i_0 < i_1$ then

$$k^{i_0n-1}(c(H)) \stackrel{\#}{\rightarrow} k^{(i_0+1)n-1}(c(H)) \stackrel{\#}{\rightarrow} \dots \stackrel{\#}{\rightarrow} k^{i_1n-1}(c(H))$$

hence

$$k^{i_0 n - 1}(c(H)) = k^{(i_0 + 1)n - 1}(c(H))$$

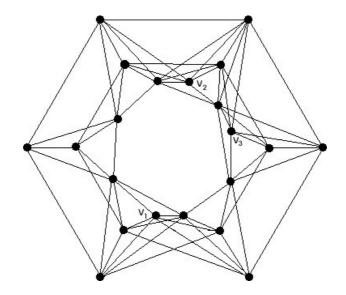
Let l be the period of c(H). Then l divide n and there is m such that $k^m(c(H)) \cong k^{m+rl}(c(H))$.

Now we prove main result of this paper:

Theorem 4.3 $G \stackrel{\#}{\to} H$ then G and H have the same k – behavior.

- 1) It is clear from the theorem 3.2.
- **2a)** If G is stationary, then there are n, m such that
- $k^n(G) \cong k^{n+m}(G) \xrightarrow{\#} k^{n+im}(H)$, for every $i \in \mathbb{N}$, but since $k^m(G)$ is dismantliable to a finite number of graphs, there are distincts i_0 and i_1 such that $k^{m+i_0n}(H) = k^{m+i_1n}(H)$. Therefore H is a stationary.
- **2b)** That H is k-stationary implies that G is k-stationary follows from Lemma 4.1.4, Lemma 4.2 and 2a).
 - 3) This is a consequence of the 1) and 2).

Remark 4.4 The graph in figure 4 has a period 3 and has three dominated points v_1, v_2, v_3 . If we delete v_1 we obtain a 6-periodic graph, if we delete v_3 we obtain a 1-periodic graph. Let G, H be graphs such that $G \stackrel{\#}{\to} H$ and there are n, and m such that $k^{m+n}(G) \cong k^m(G)$. From Theorem 4.3 there are m' and n' such that $k^{m'+n'}(H) \cong k^{m'}(H)$. The determination of the relation among m, n, m' and n' is an open problem.



References

- [1] F. Escalante. Über Iterierte Clique-Graphen. Abh. Math. Sem. Univ. Hamburg. 39 (1973) 58-68.
- [2] M.E. Frías-Armenta. Tesis Doctoral: *Gráficas Iteradas de Clanes*. Facultad de Ciencias, Universidad Nacional Autonoma de México (2000) Dir. V. Neumann-Lara Imunam.
- [3] M.E. Frías-Armenta, V. Neumann-Lara. Cliques, Prekernels, Operators, Covers and pointer graphs. In preparation.
- [4] F. Harary. Graph Theory. Addison-Wesley, Reading, MA (1969).
- [5] S. Hazan, V. Neumann-Lara. Fixed Points of Posets and Clique Graph. Kliwer Academic Publishers 13 (1996) 219-225.
- [6] V. Neumann-Lara. On Clique-divergent Graphs, Problèmes Combinatoires et Théorie des Graphes (Colloques internationaux C.N.R.S, 260). Paris (1978), 313-315.
- [7] V. Neumann-Lara. Clique Divergence in Graphs, Algebraic Methods in Graph Theory (Coll. Math.Soc.Janos Bolyai, 25). Szeged (1981), 563-569.
- [8] E. Prisner. Convergence of Iterated Clique Graphs. Discrete Mathematics 103 (1992) 199-207.